

Last-Level Cache Side-Channel Attacks Are Feasible in the Modern Public Cloud

Zirui Neil Zhao, Adam Morrison, Christopher W. Fletcher, Josep Torrellas

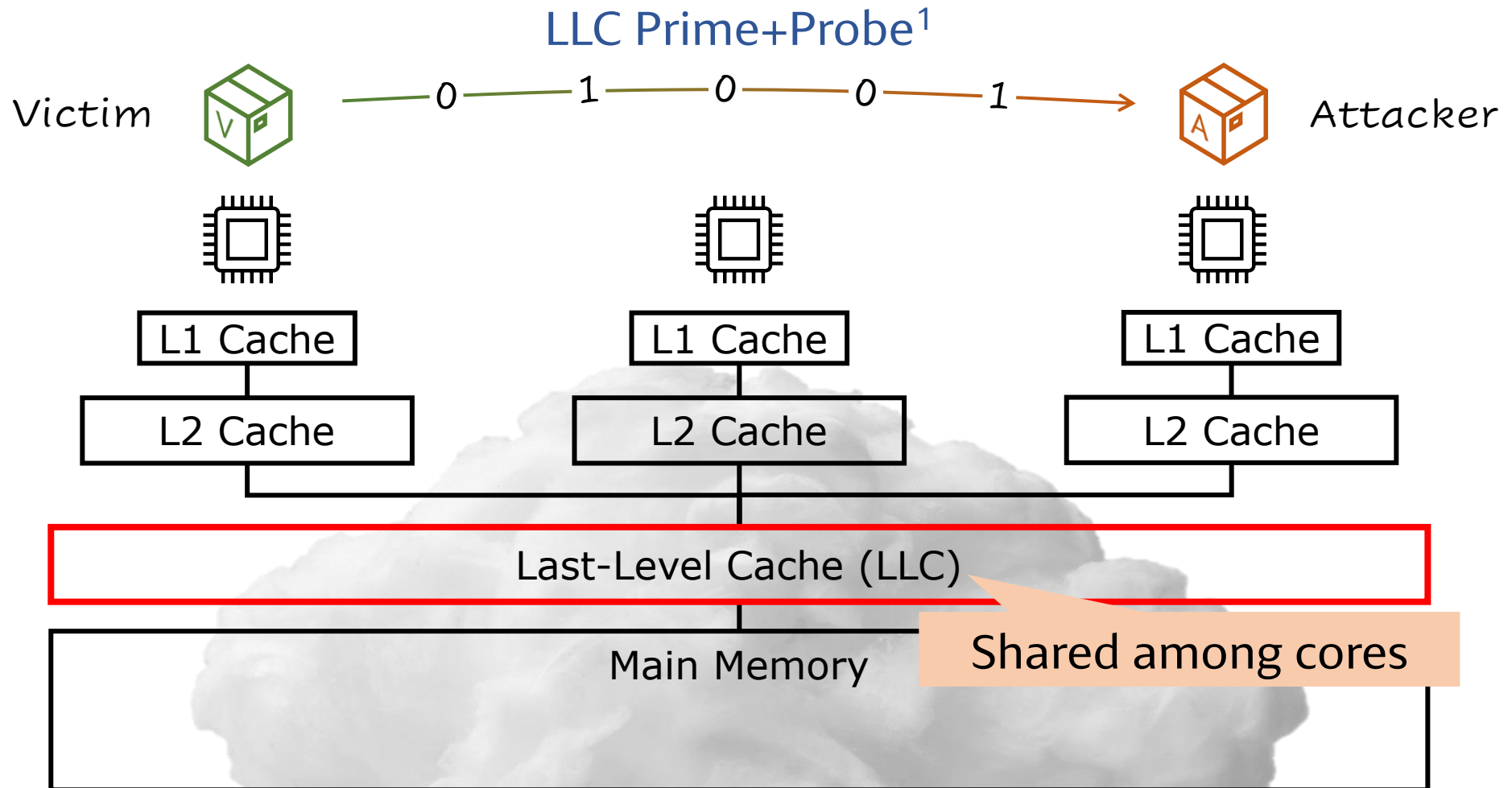
University of Illinois Tel Aviv University

ASPLOS '24 – Session 2B: Side Channels

I ILLINOIS



Shared Last-Level Cache (LLC) Enables Information Leakage



¹Liu et al., Last-Level Cache Side-Channel Attacks are Practical (S&P 2015)

Cloud Vendors Claim LLC Prime+Probe is Impractical

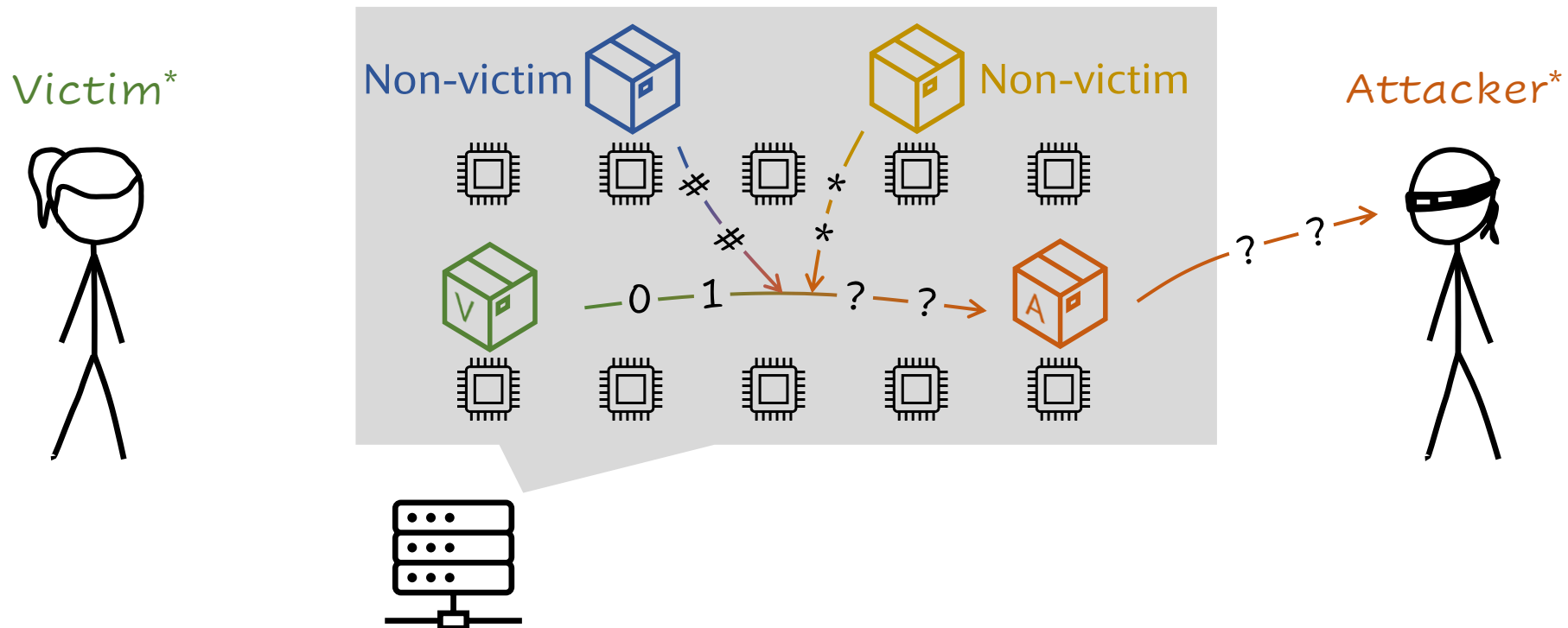
Example: AWS Whitepaper – The Security Design of the AWS Nitro System
(Version November 18, 2022)



Paraphrased:

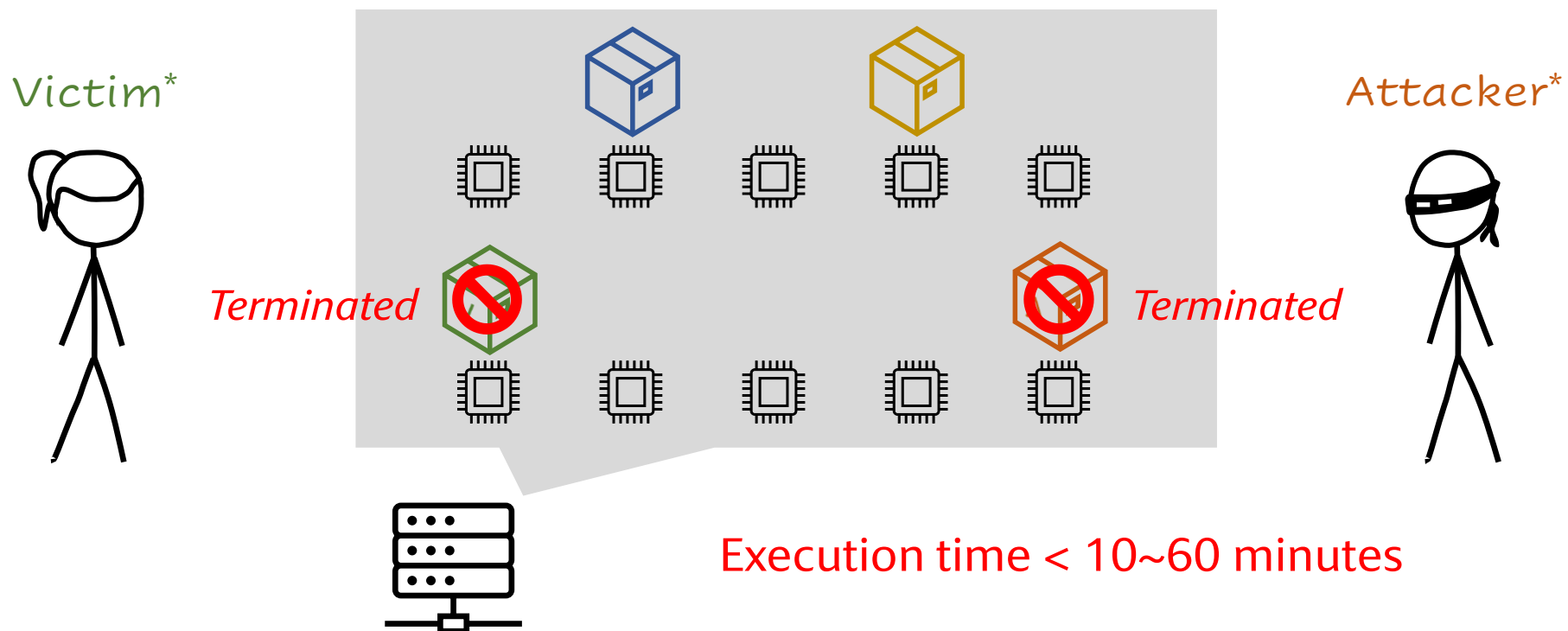
Last-level cache (LLC) Prime+Probe is impractical due to the noise; therefore, our side-channel mitigations are **very strong** even if we do not protect VMs against LLC Prime+Probe

Attacker's Challenge 1: Production Cloud is Noisy



*Characters are based on <https://xkcd.com/2176> and <https://xkcd.com/1808> (under a CC Attribution-NonCommercial 2.5 License)

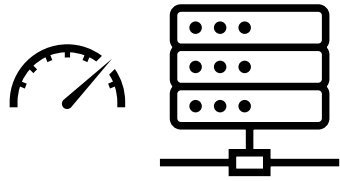
Attacker's Challenge 2: Modern Clouds (e.g., FaaS) are Dynamic



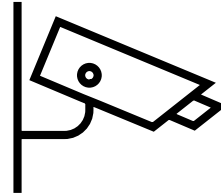
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Contributions of This Work (at 31,000 Feet)

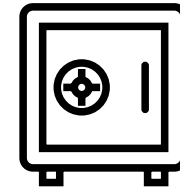
Cross-tenant information leakage with LLC Prime+Probe  Google Cloud



1. Fast LLC Channel Setup
> 10 hours → 2.4 minutes



2. Noise-Resilient Victim
Monitoring

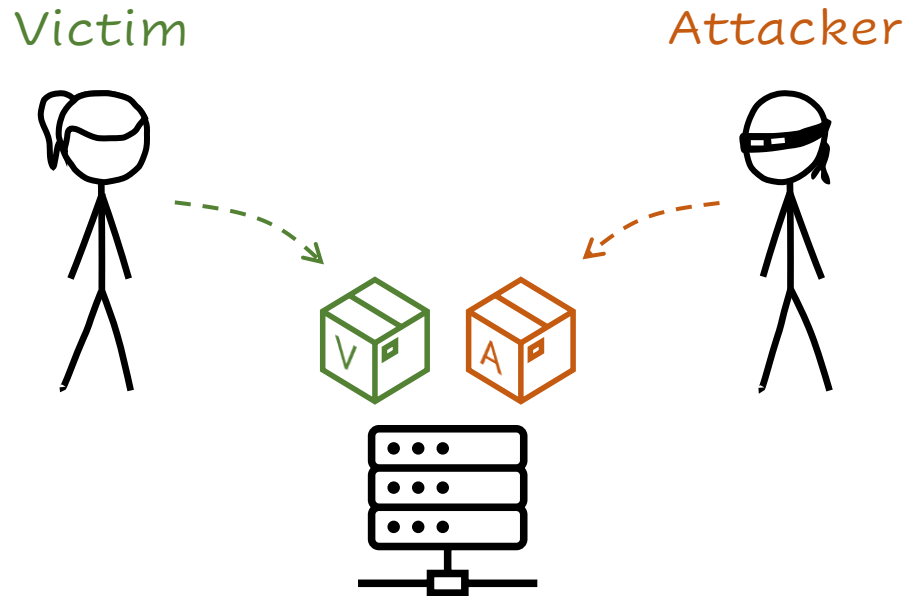


3. Information Extraction

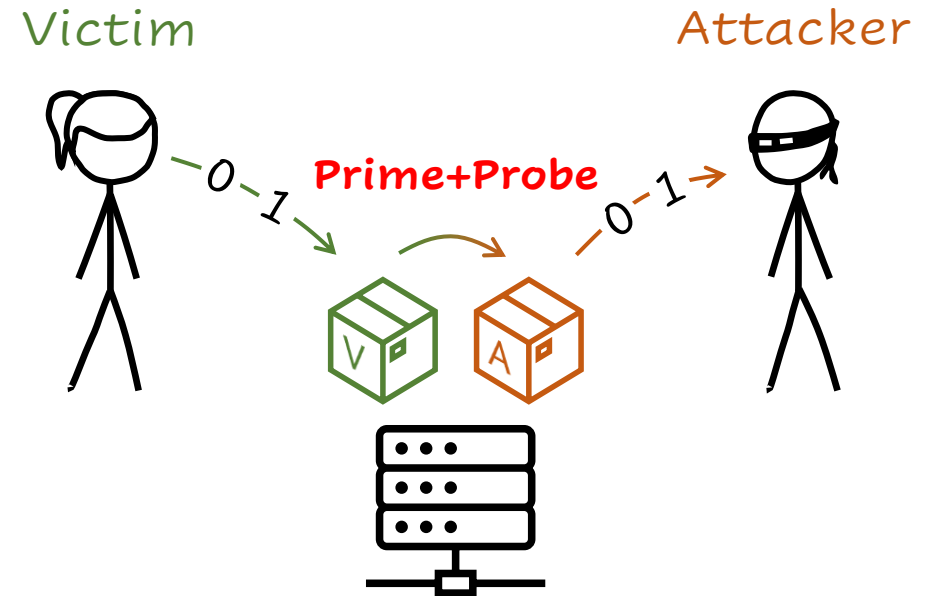
- + Google filed a **critical-level bug** to their product team
- + AWS revised their whitepaper on February 15, 2024

Threat Model & Assumptions

✓ Step 1: Co-Location



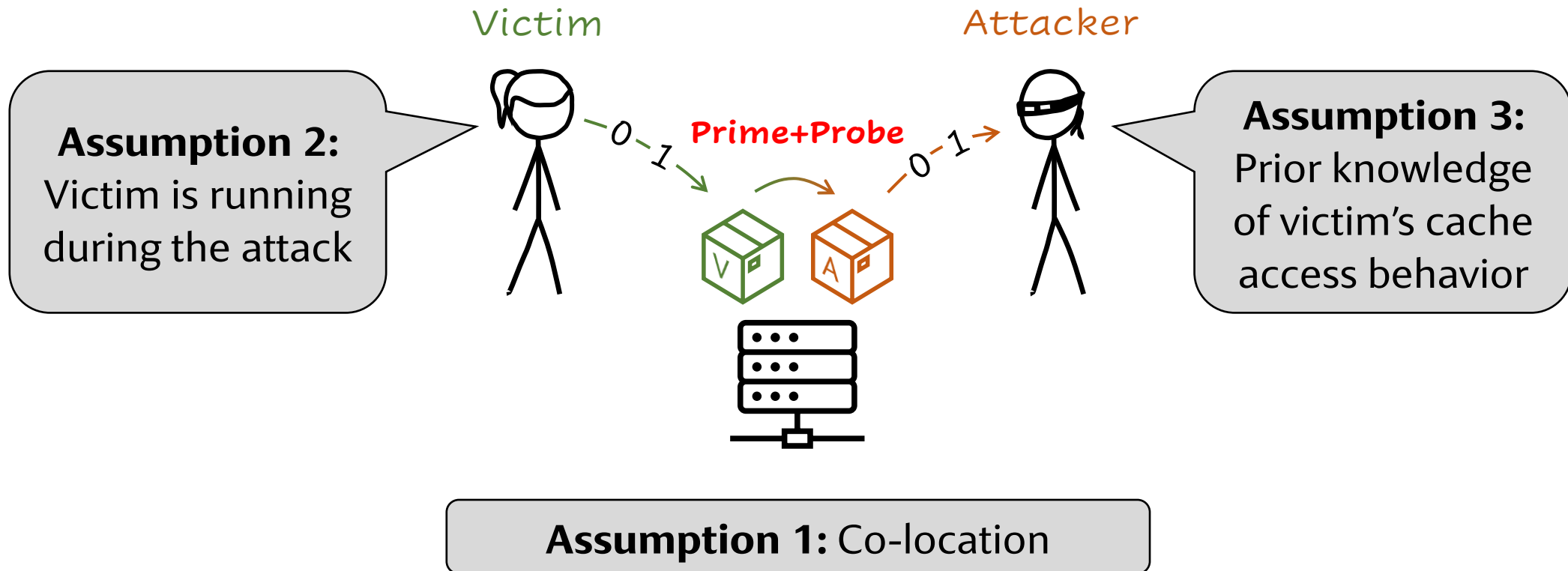
Step 2: LLC Prime+Probe



✓ Everywhere All at Once: Co-Location
Attacks on Public Cloud FaaS
(ASPLOS '24 – Session 1D)

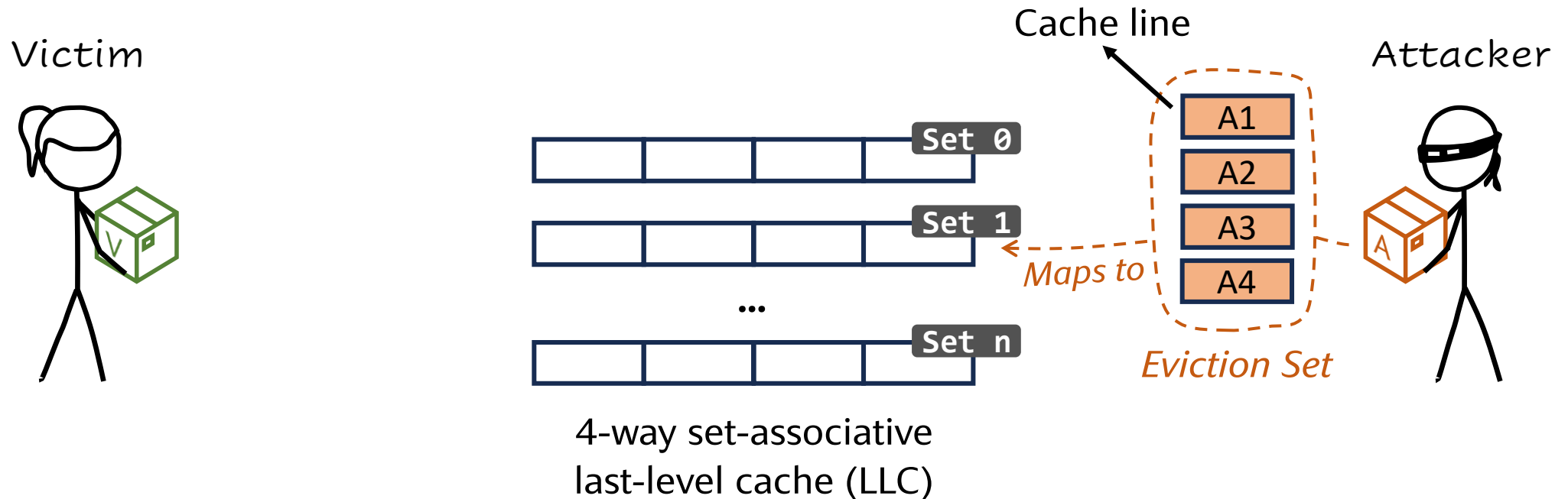
Threat Model & Assumptions

Step 2: LLC Prime+Probe



Background: LLC Prime+Probe Attack

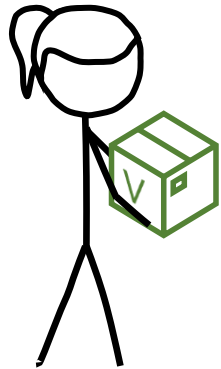
Eviction set: A set of cache lines that fully occupy a cache set



Background: LLC Prime+Probe Attack

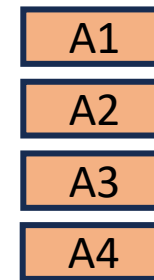
Eviction set \Rightarrow Monitor memory accesses to an LLC set with **Prime+Probe**

Victim

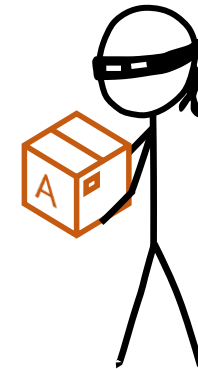


4-way set-associative
last-level cache (LLC)

Attacker



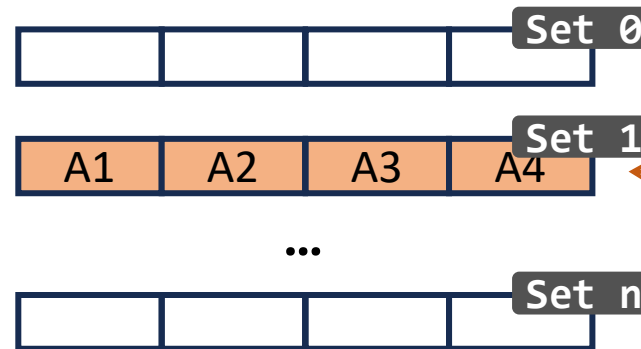
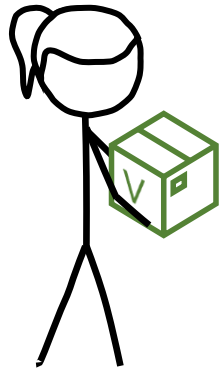
Eviction Set



Background: LLC Prime+Probe Attack

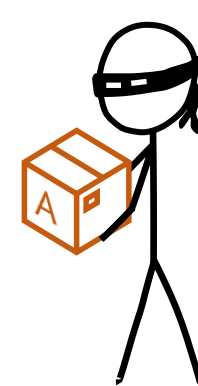
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last-level cache (LLC)

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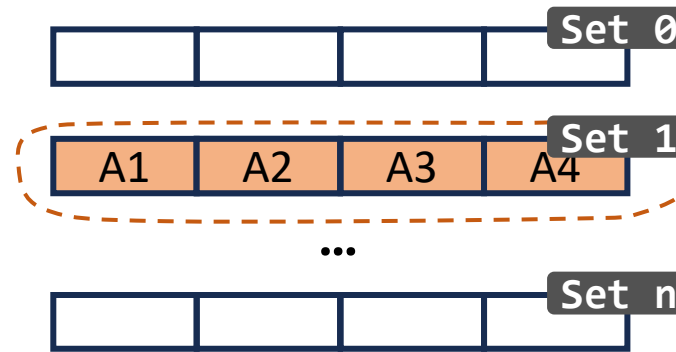
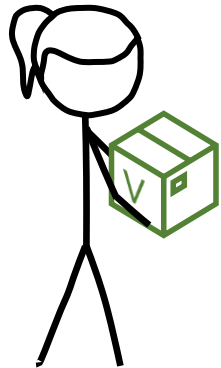


Eviction Set

Background: LLC Prime+Probe Attack

Eviction set \Rightarrow Monitor memory accesses to an LLC set with **Prime+Probe**

Victim

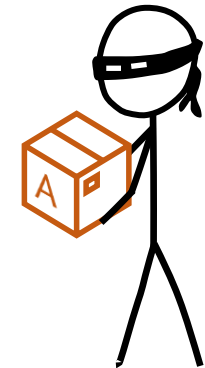


4-way set-associative
last-level cache (LLC)

Probe

Cache hit!
Low latency

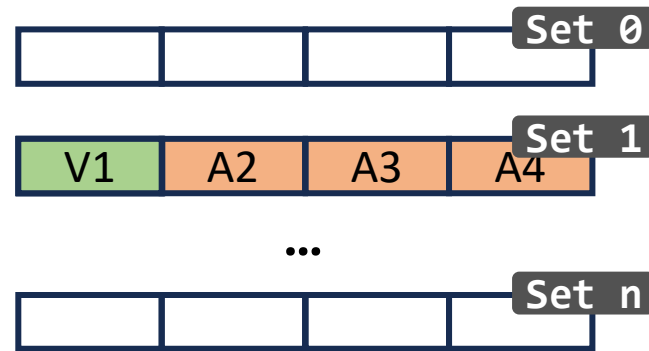
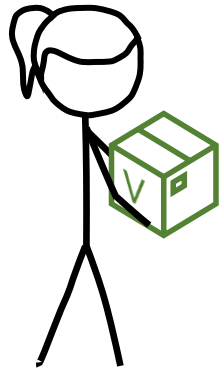
Attacker



Background: LLC Prime+Probe Attack

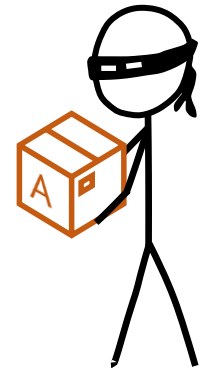
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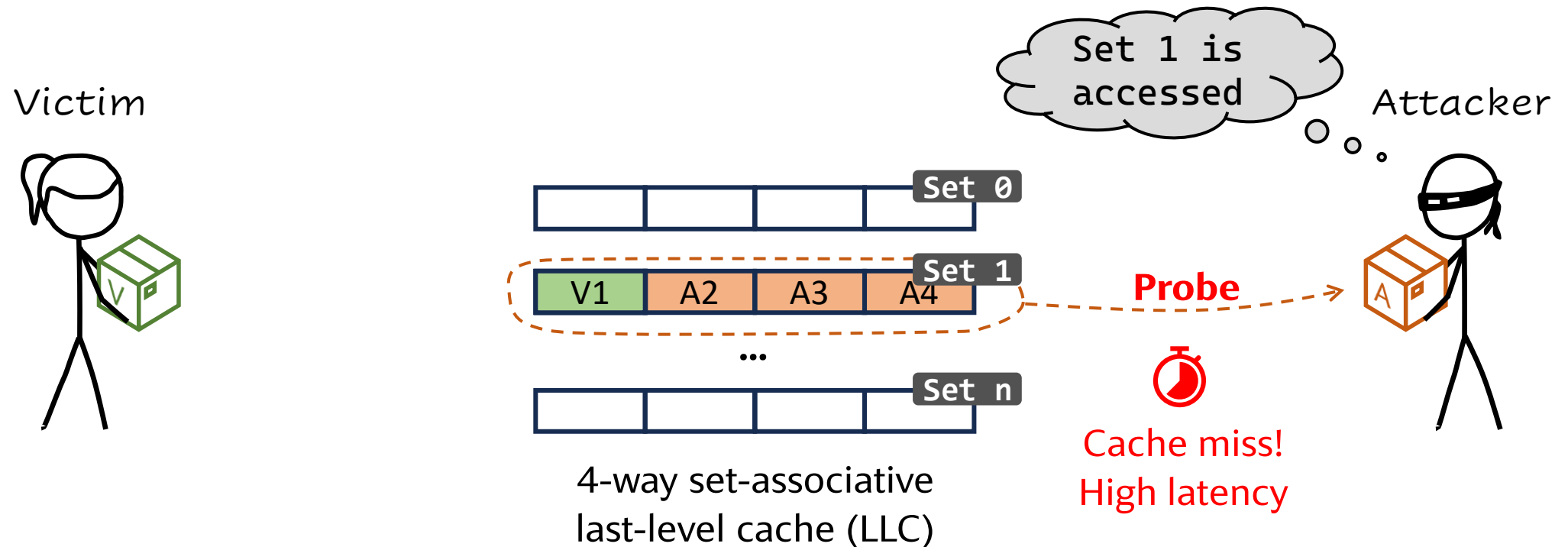
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Attacker



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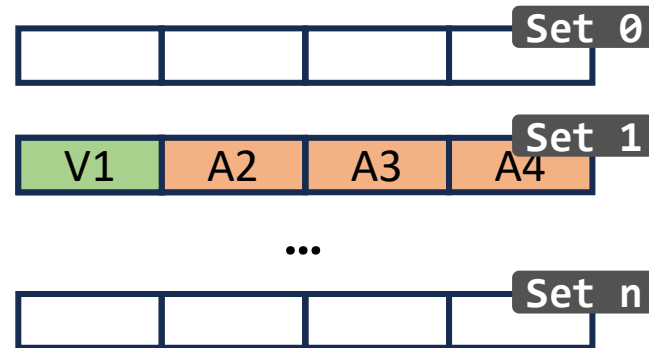
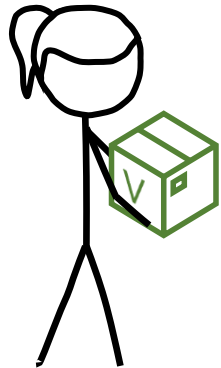


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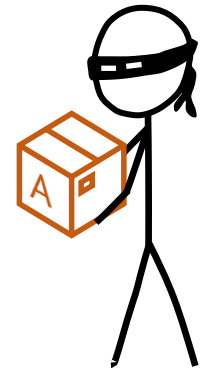
Takeaway

Victim



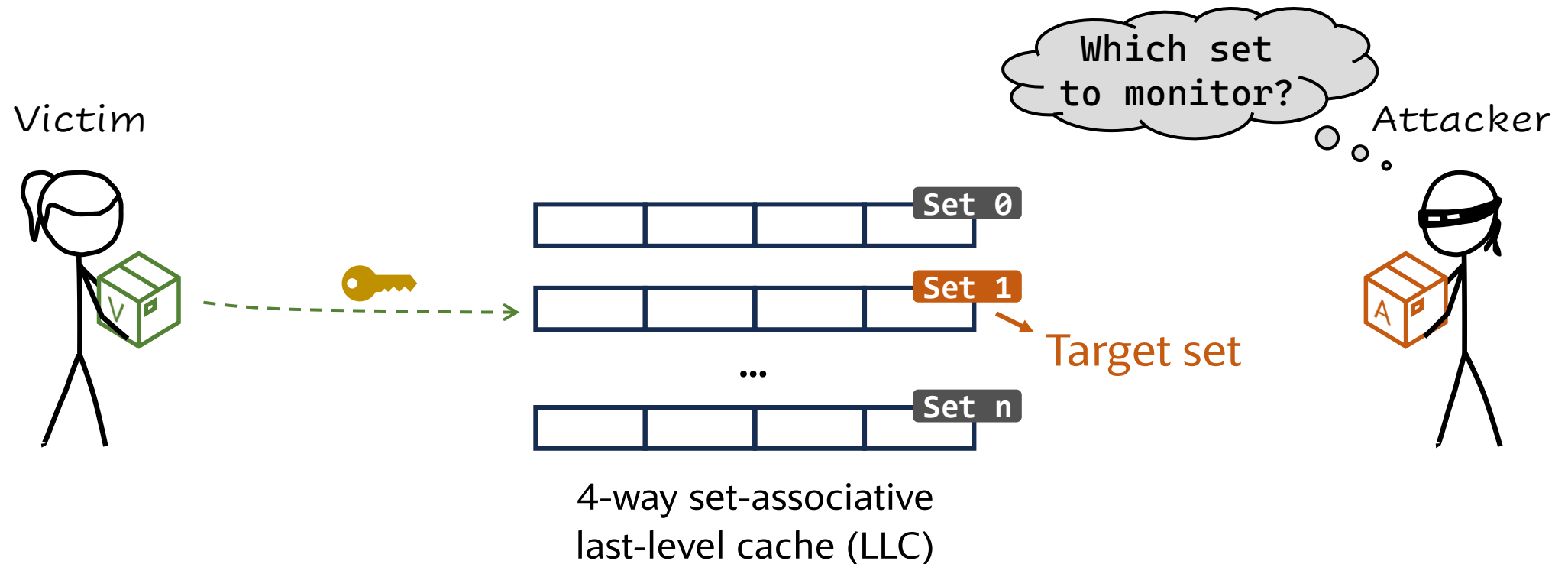
4-way set-associative
last-level cache (LLC)

Attacker



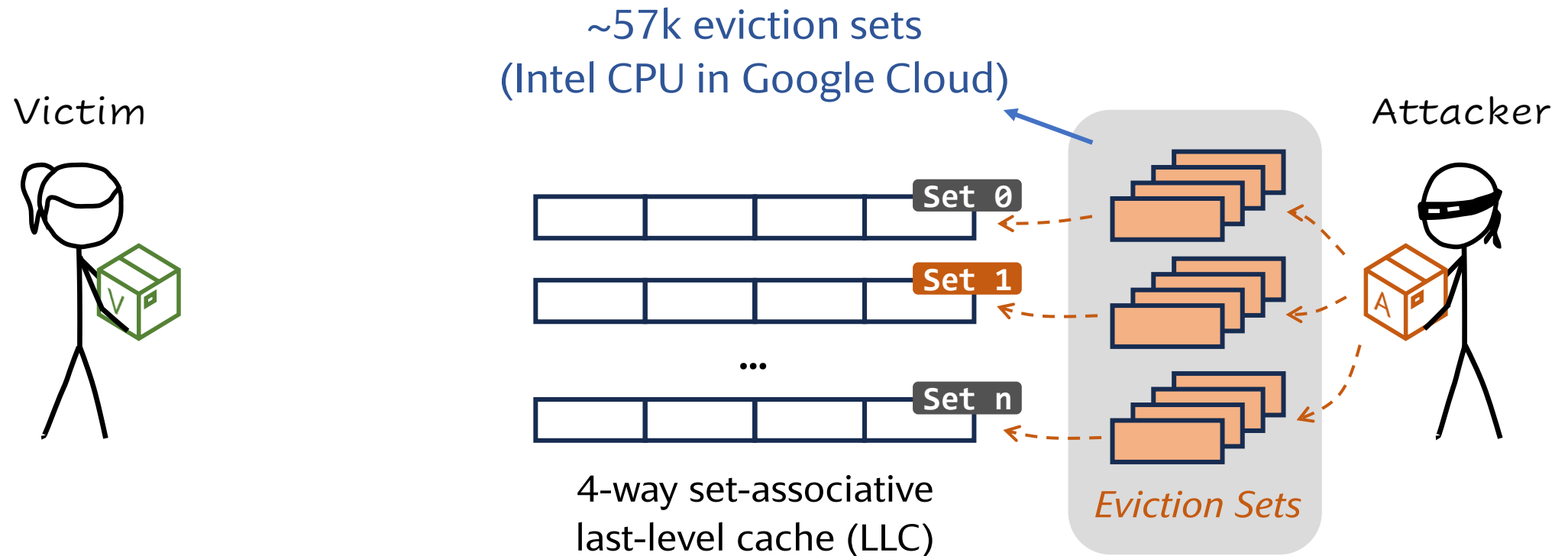
An Unprivileged Attacker Does Not Know the Target Set

Target set: An LLC set accessed by the victim in a secret-dependent manner



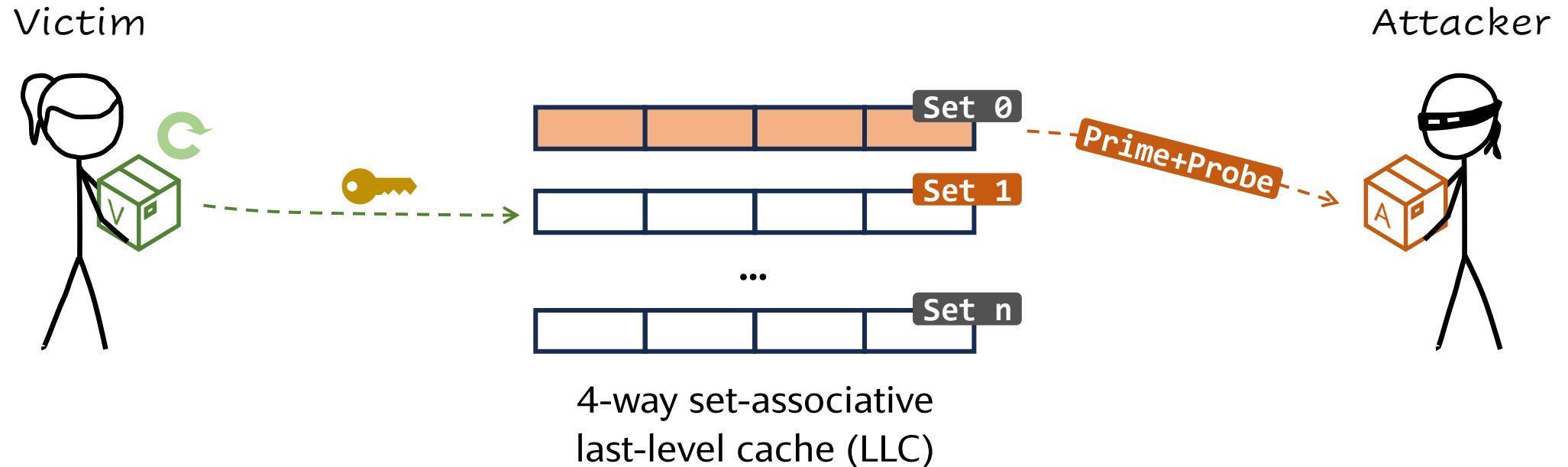
Step 2.1: Build Many Eviction Sets

Attacker needs an eviction set for every LLC set in the system



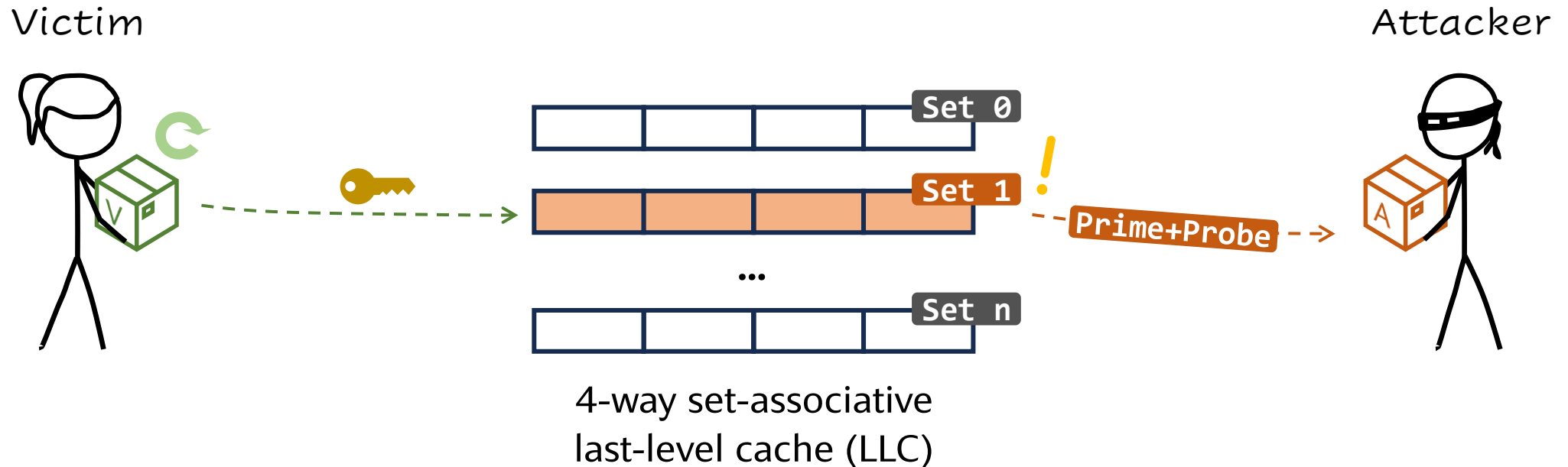
Step 2.2: Identify Target LLC Set to Monitor

Attacker collects an access trace from *each* LLC set
⇒ Checks whether the access trace matches victim's access behavior



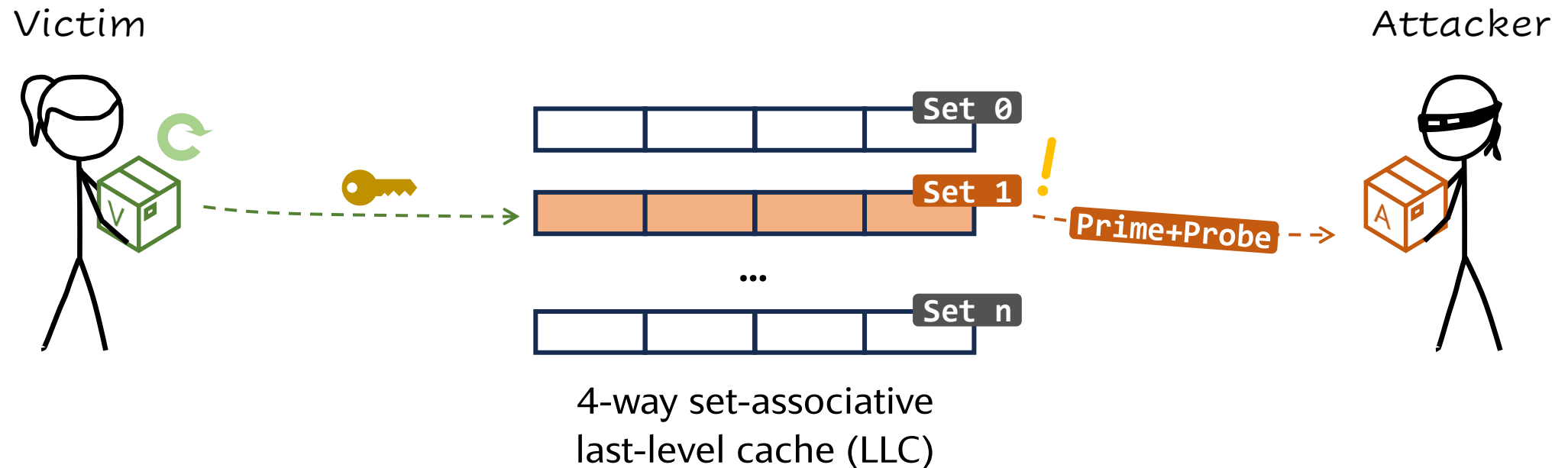
Step 2.2: Identify Target LLC Set to Monitor

Attacker collects an access trace from *each* LLC set
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Step 2.3: Extract Information from the Victim

Attacker monitors the target set and extracts the sensitive information



Attack Roadmap


✓ Step 1: Co-Locate *— Zhao et al., Everywhere All at Once... (ASPLOS '24)*

This Work — Step 2.1: Build Many Eviction Sets
— Step 2.2: Identify Target Set
— Step 2.3: Extract Information

} Challenges: Noise and dynamism

↪ End-to-end, cross-tenant information leakage in production Google Cloud

Attack Roadmap

✓ Step 1: Co-Locate  *Zhao et al., Everywhere All at Once... (ASPLOS '24)*

➔ **Step 2.1: Build Many Eviction Sets**

Step 2.2: Identify Target Set

Step 2.3: Extract Information

Challenges: Noise and dynamism

End-to-end, cross-tenant information leakage in production Google Cloud

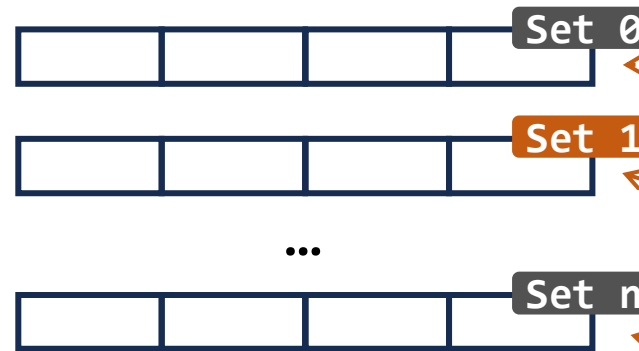
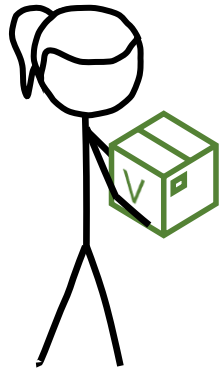
Our Contribution: Fast Eviction Set Construction

Contribution: Fast eviction set construction in *~2.4 minutes*  Google Cloud

Existing approaches: would take *>10 hours* to complete due to noise

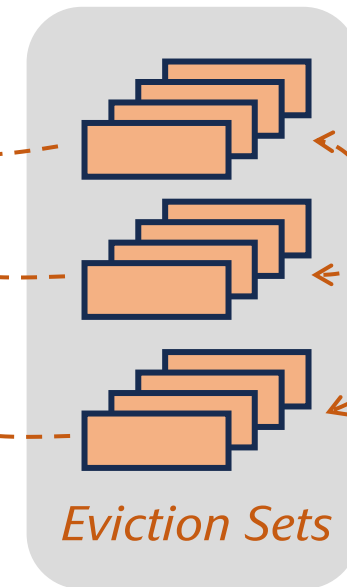
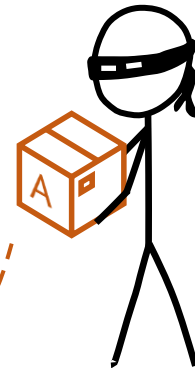
↳ Exceeds execution time limits!

Victim

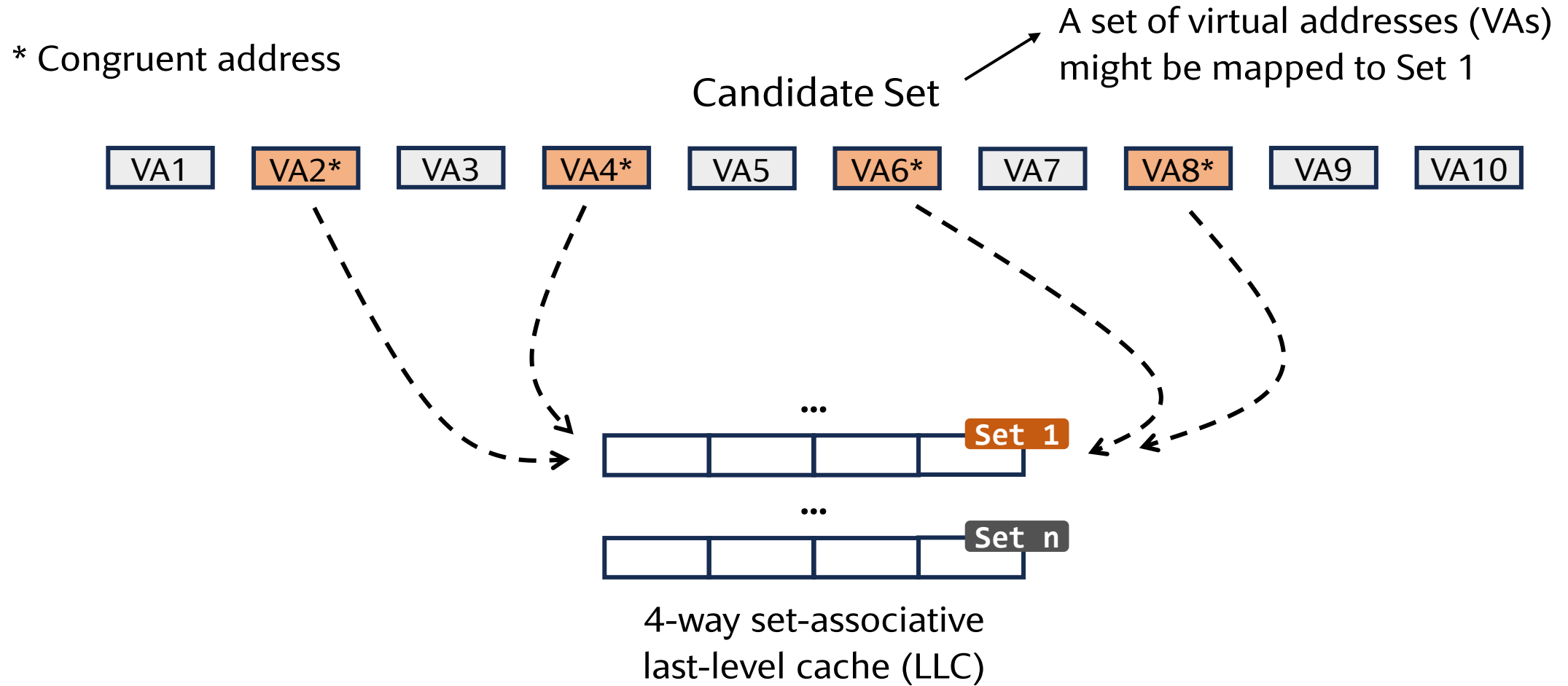


4-way set-associative
last-level cache (LLC)

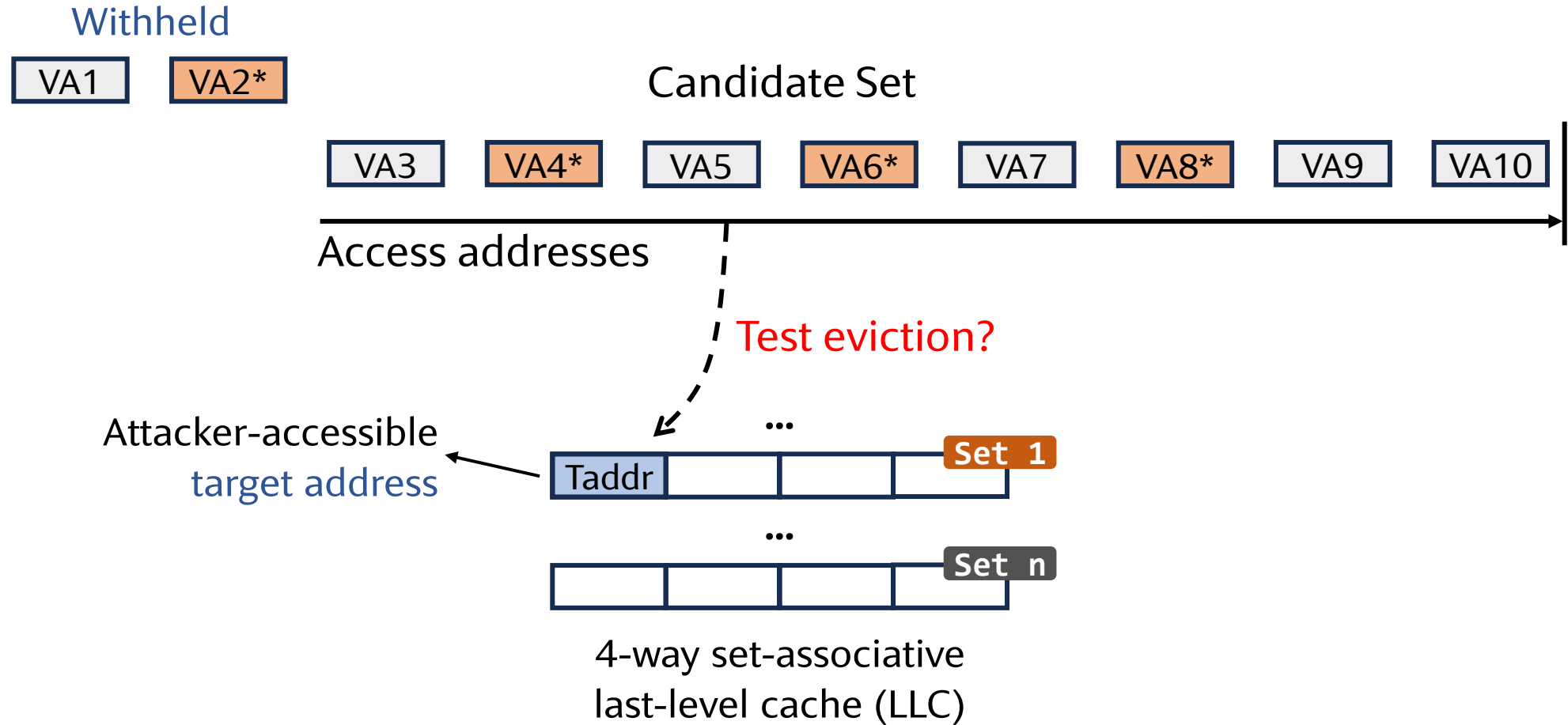
Attacker



Background: Constructing an Eviction Set



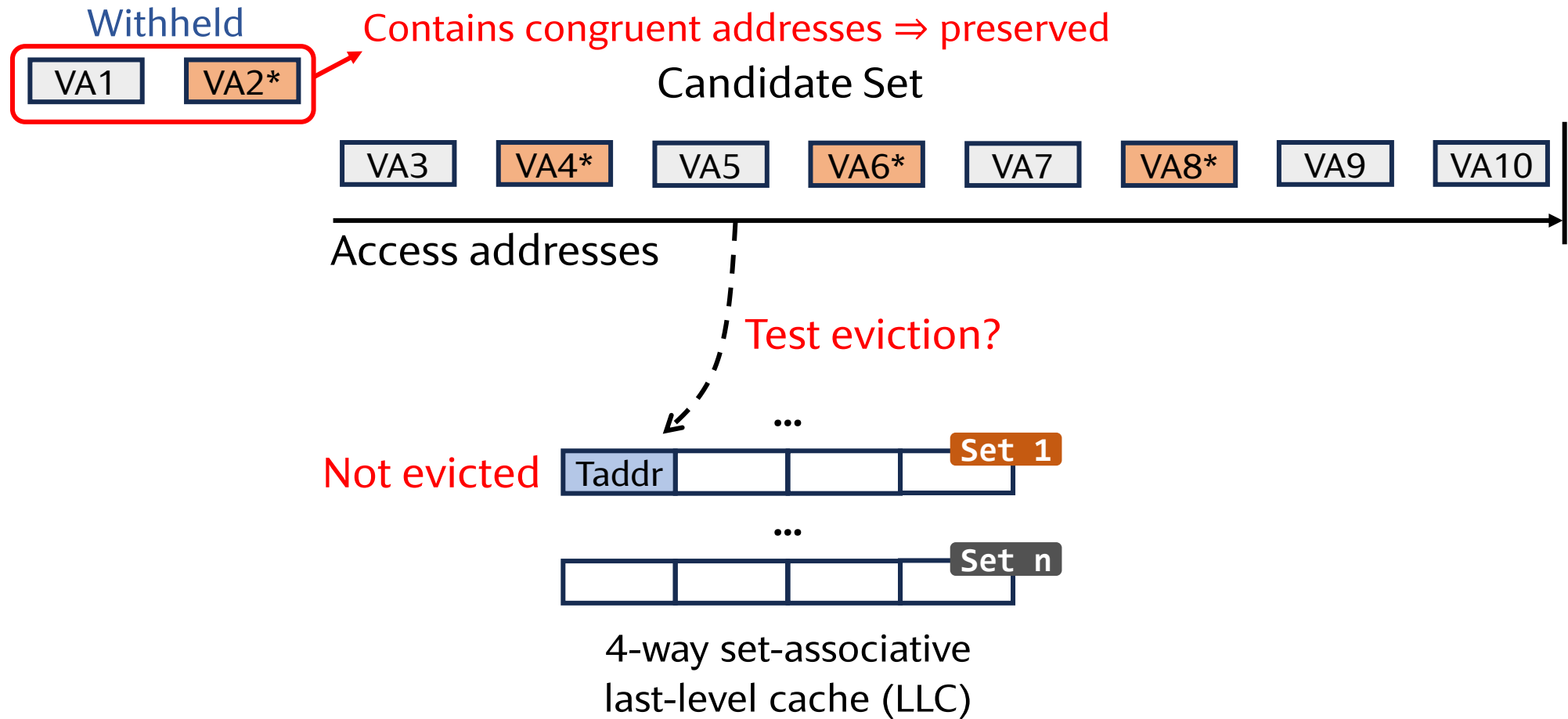
Background: Pruning the Candidate Set (Group Testing^{1,2})



¹Vila et al., Theory and Practice of Finding Eviction Sets (S&P 2019)

²Qureshi et al., New Attacks and Defense for Encrypted-Address Cache (ISCA 2019)

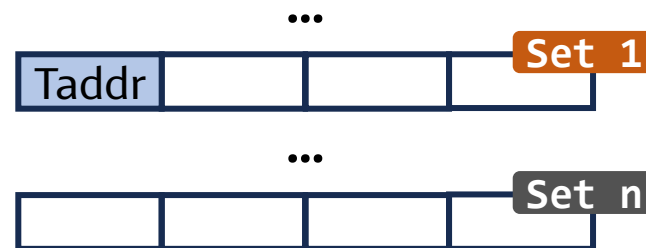
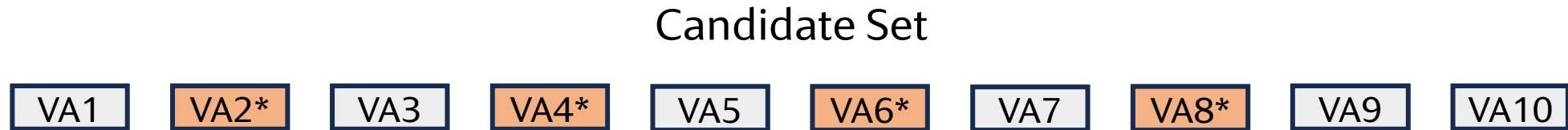
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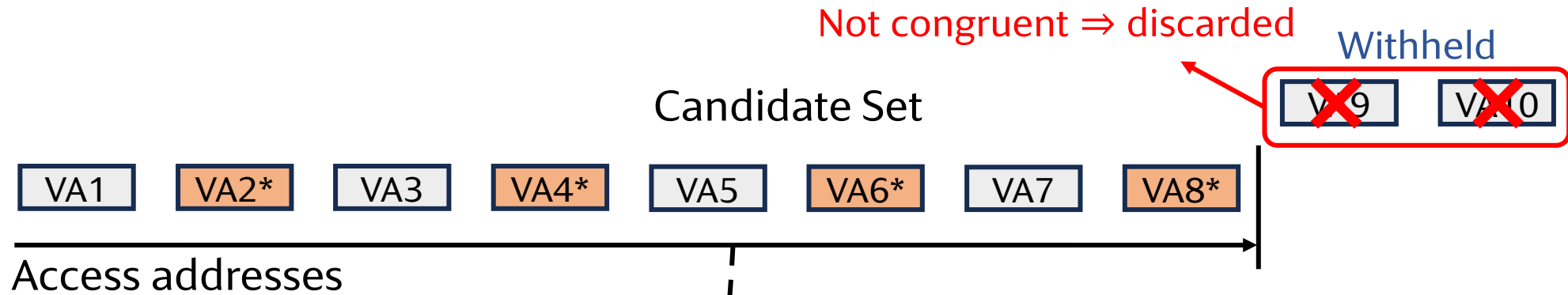


4-way set-associative
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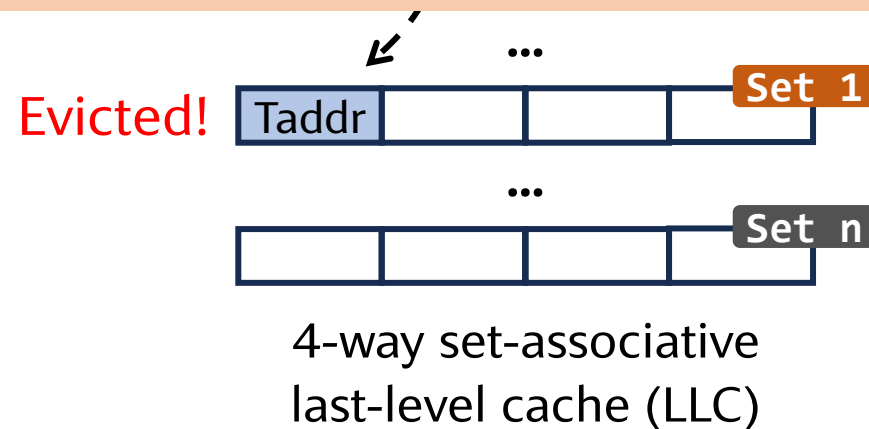
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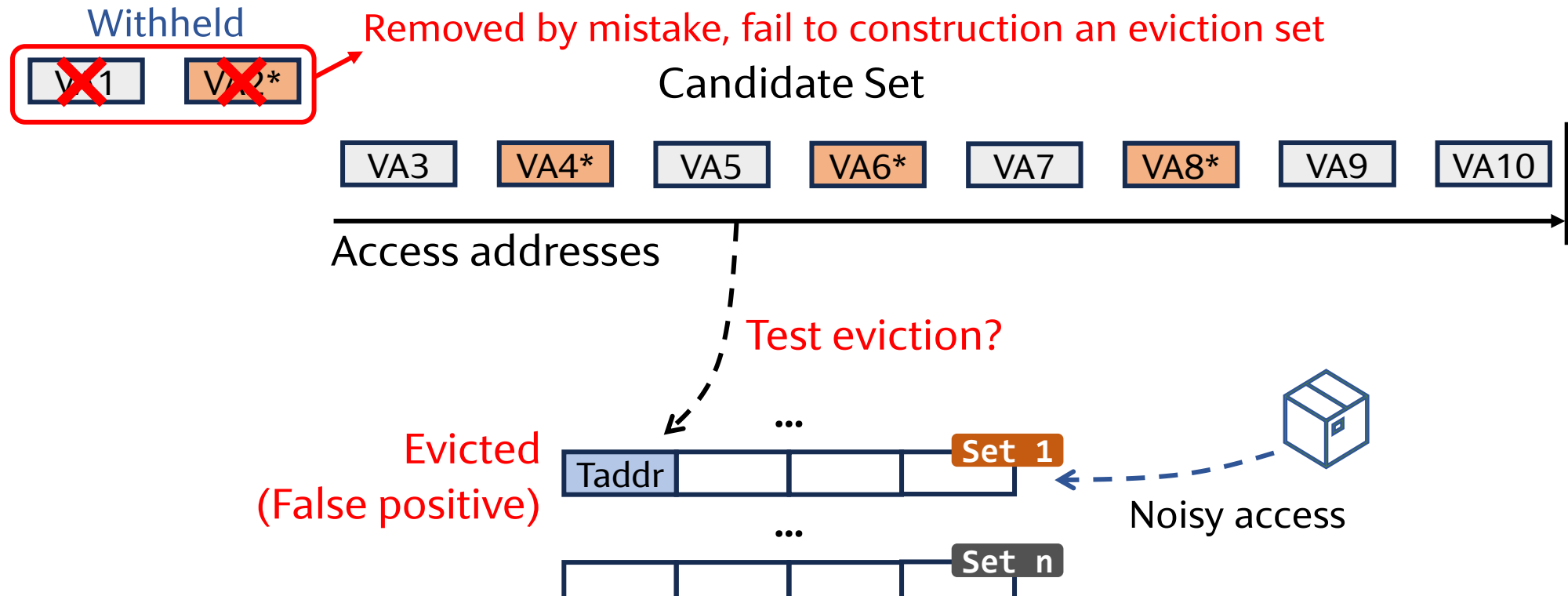
Requires $O(W^2N)$ memory accesses¹, N is the candidate set size, W is the associativity



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²Qureshi et al., New Attacks and Defense for Encrypted-Address Cache (ISCA 2019)

This Work: Test Eviction is Susceptible to Noise



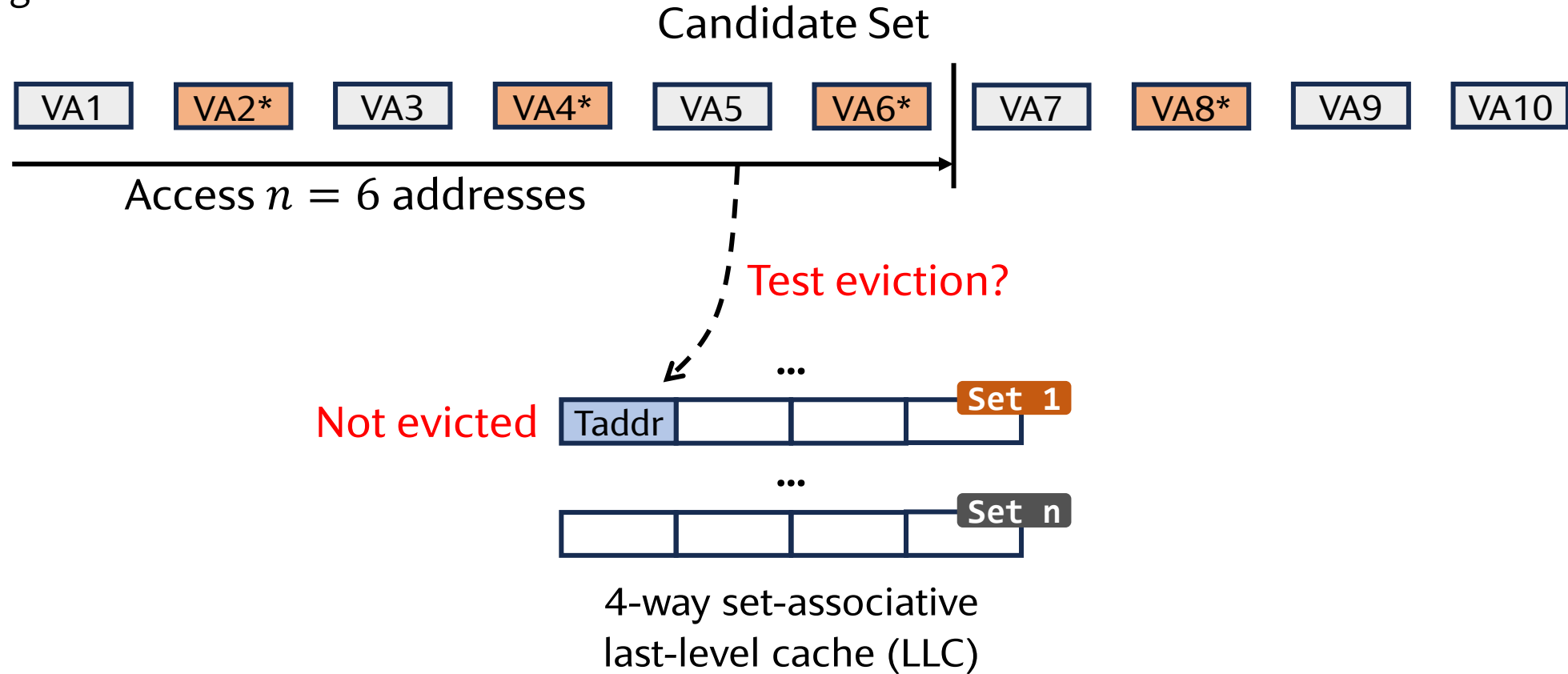
Prime+Scope¹ is similarly susceptible to noise

Our paper provides more detailed quantitative analyses of both algorithms

¹Purnal et al., Prime+Scope: Overcoming the Observer Effect for High-Precision Cache Contention Attacks (CCS 2021)

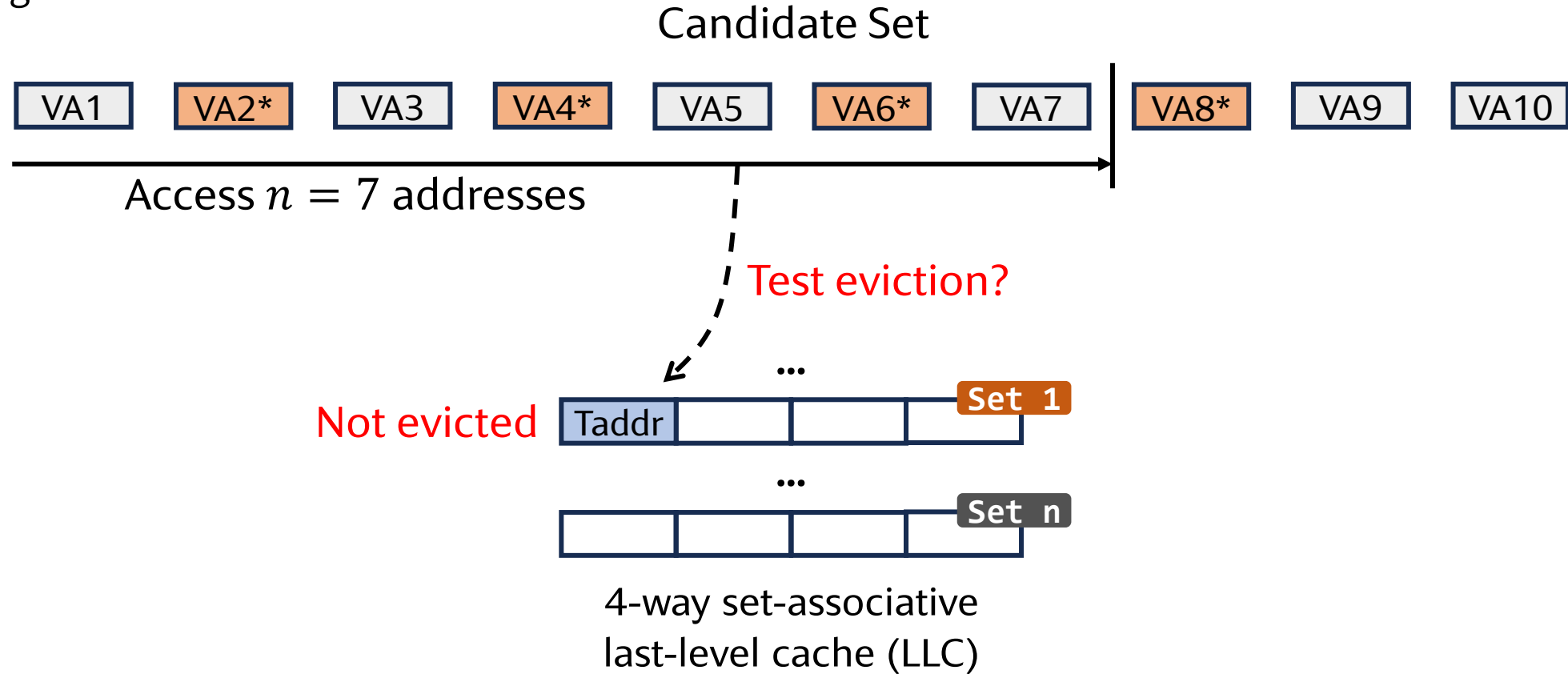
Intuition of the New Algorithm

* Congruent address

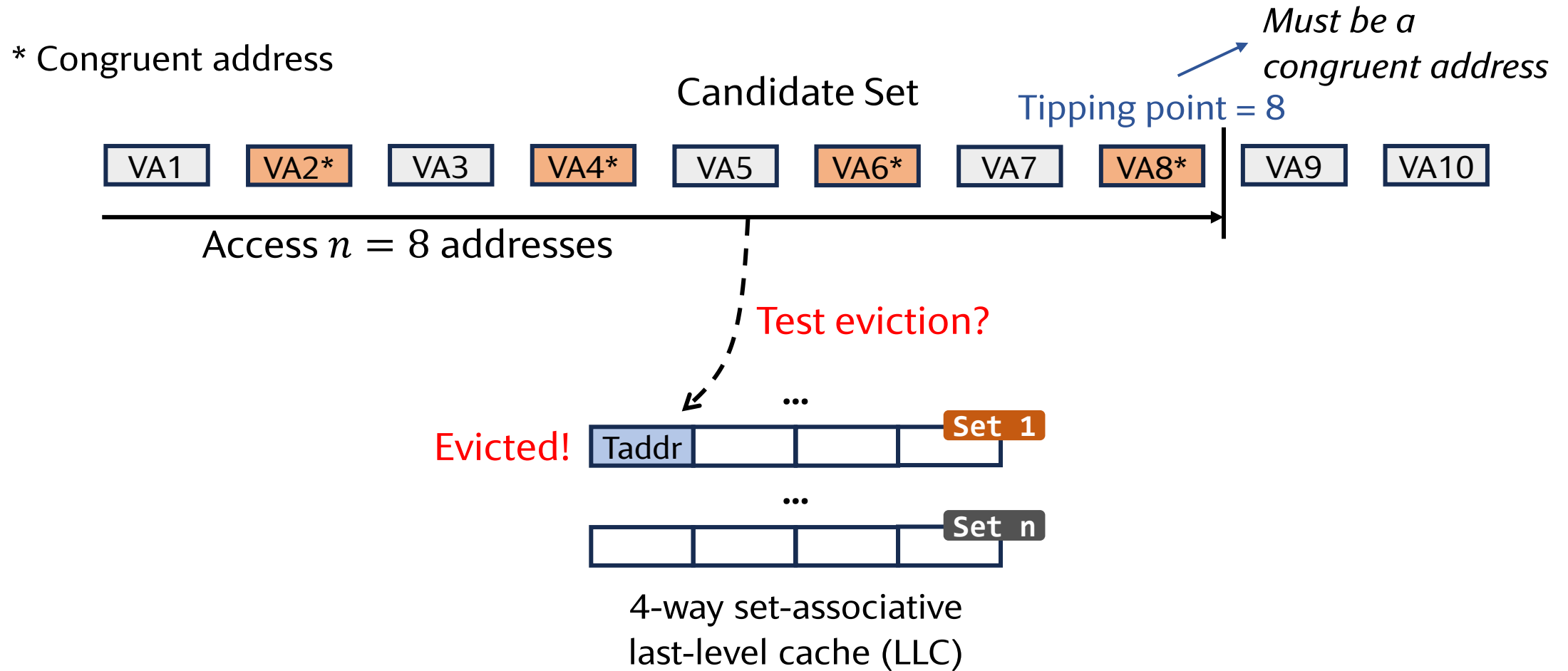


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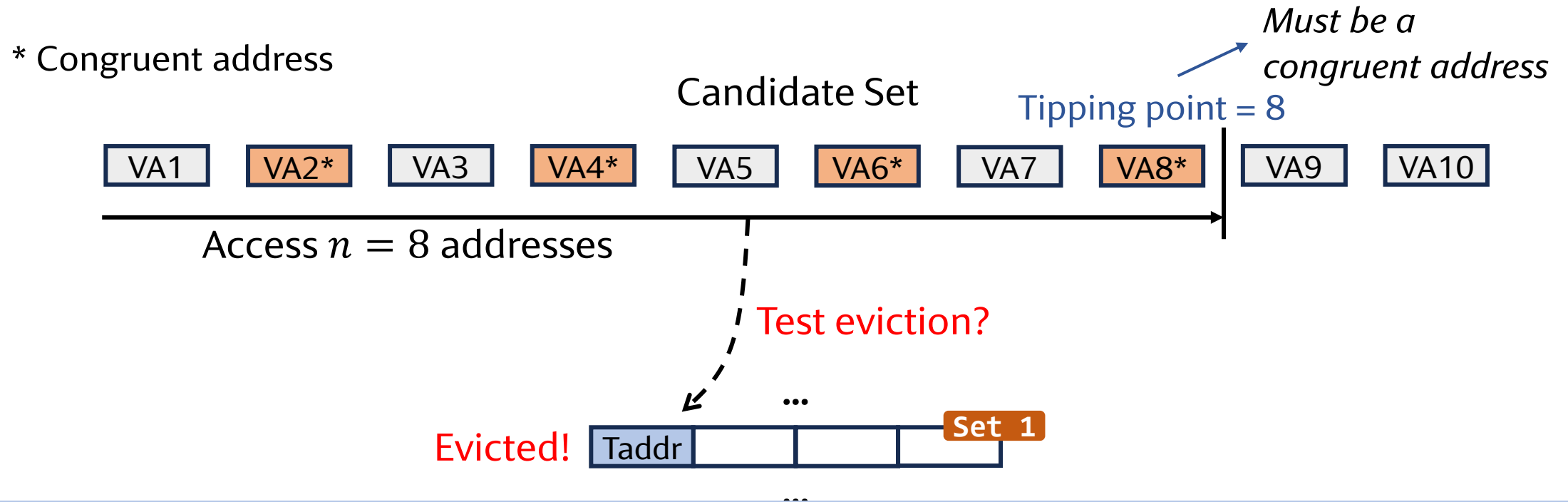
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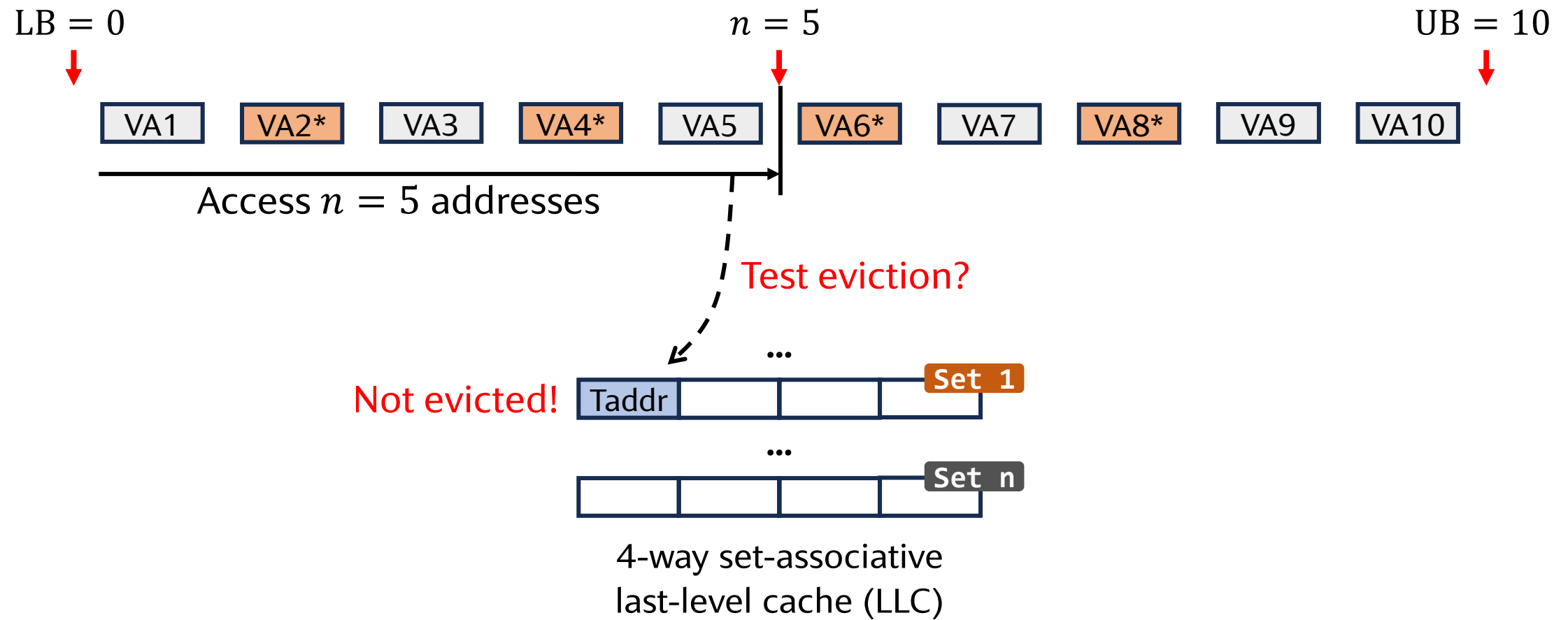


Idea: Identifying congruent addresses by finding tipping points

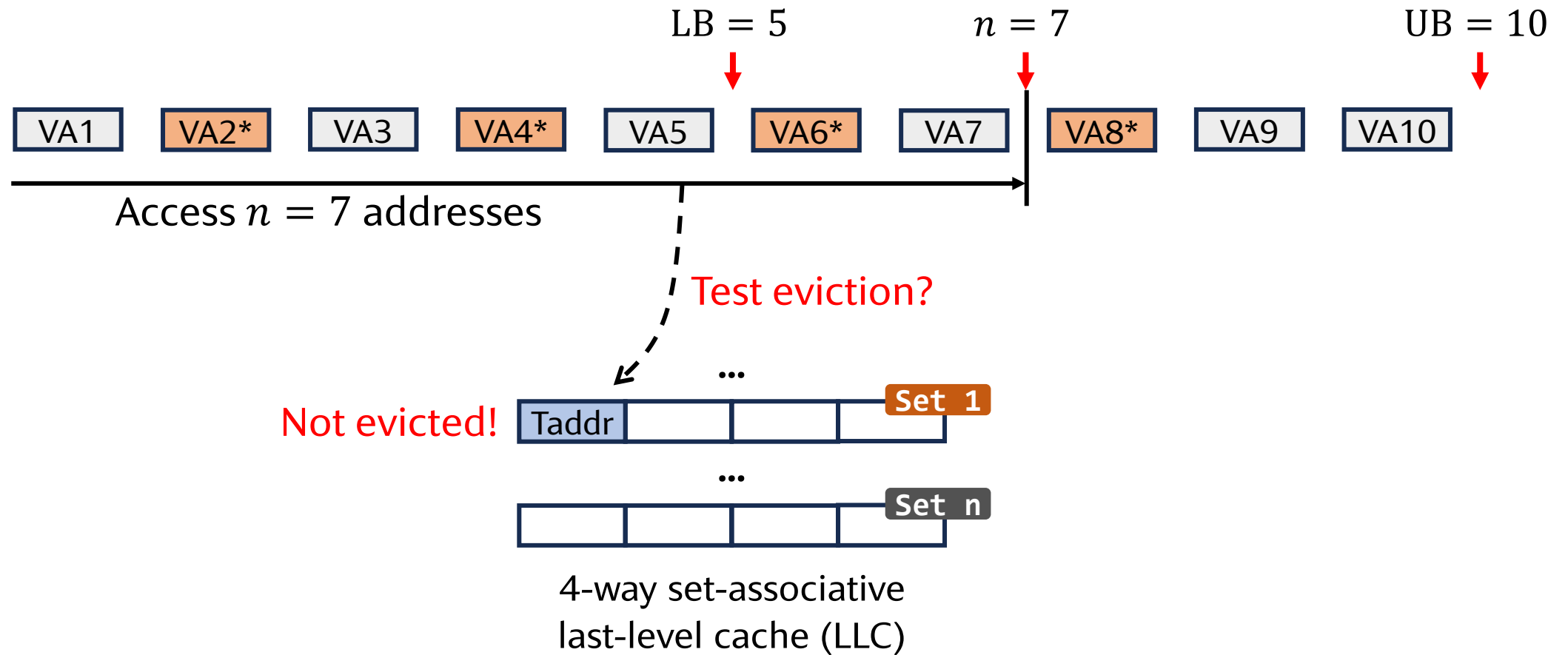
Insight: Can evict $\Rightarrow n \geq$ tipping point; Cannot evict $\Rightarrow n <$ tipping point

\Rightarrow Speed up the process with **binary search**

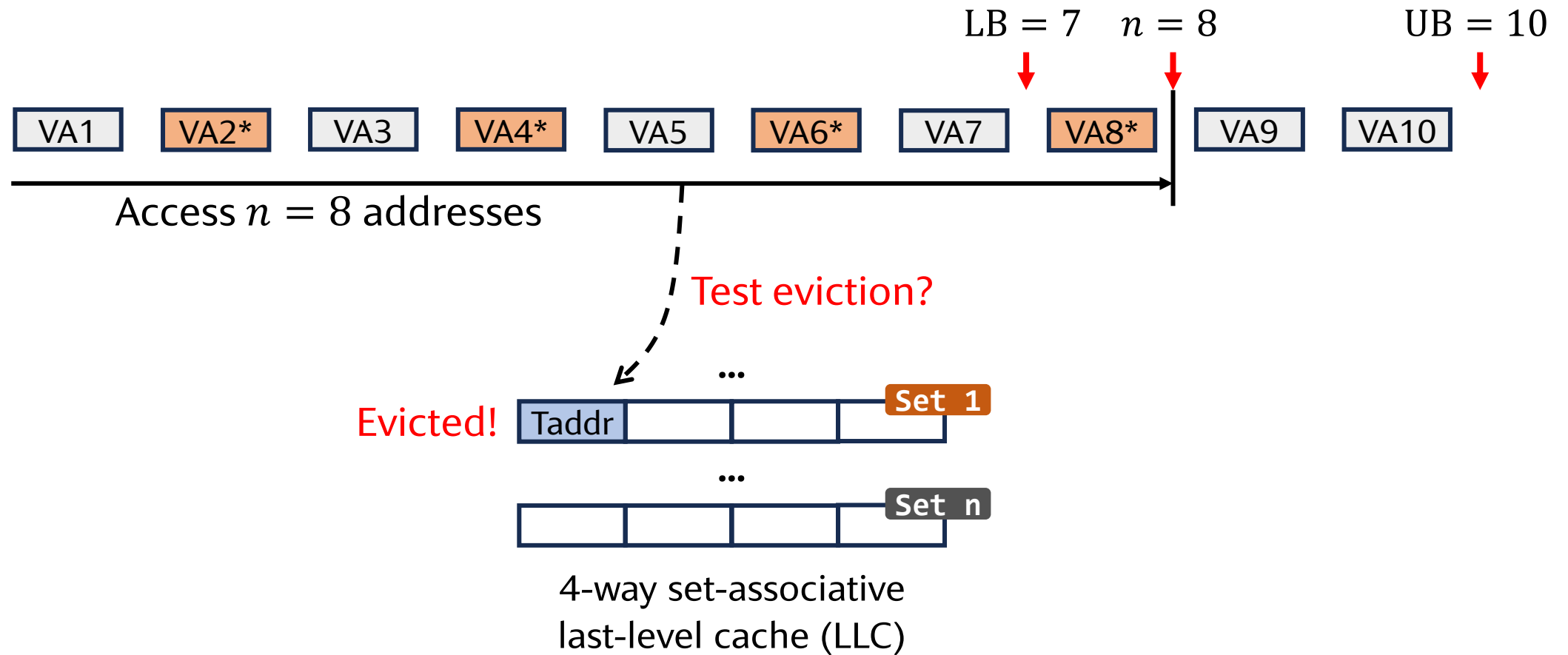
New Technique 1: A Binary Search-Based Algorithm



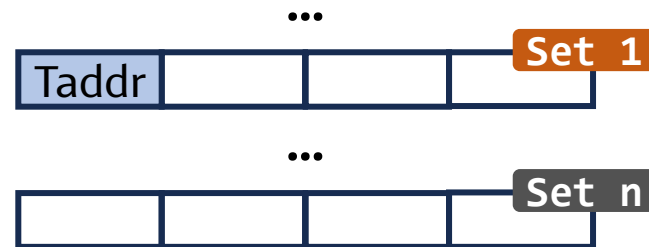
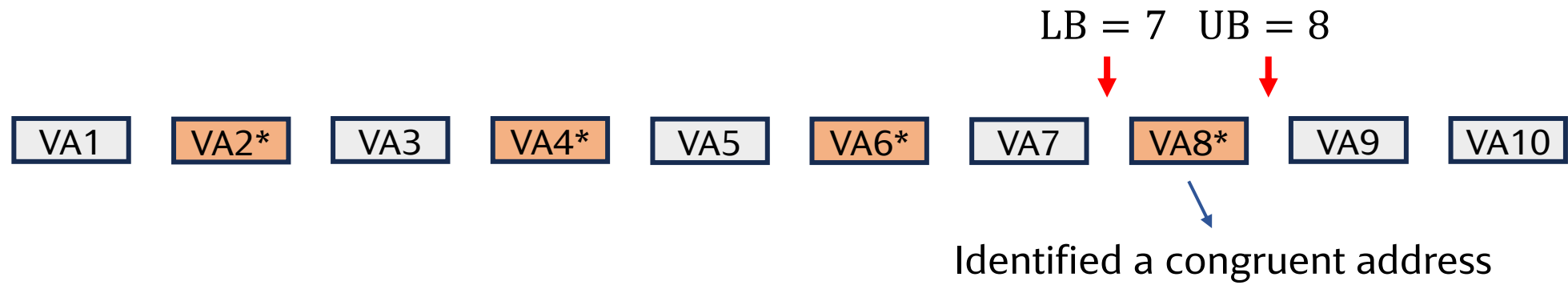
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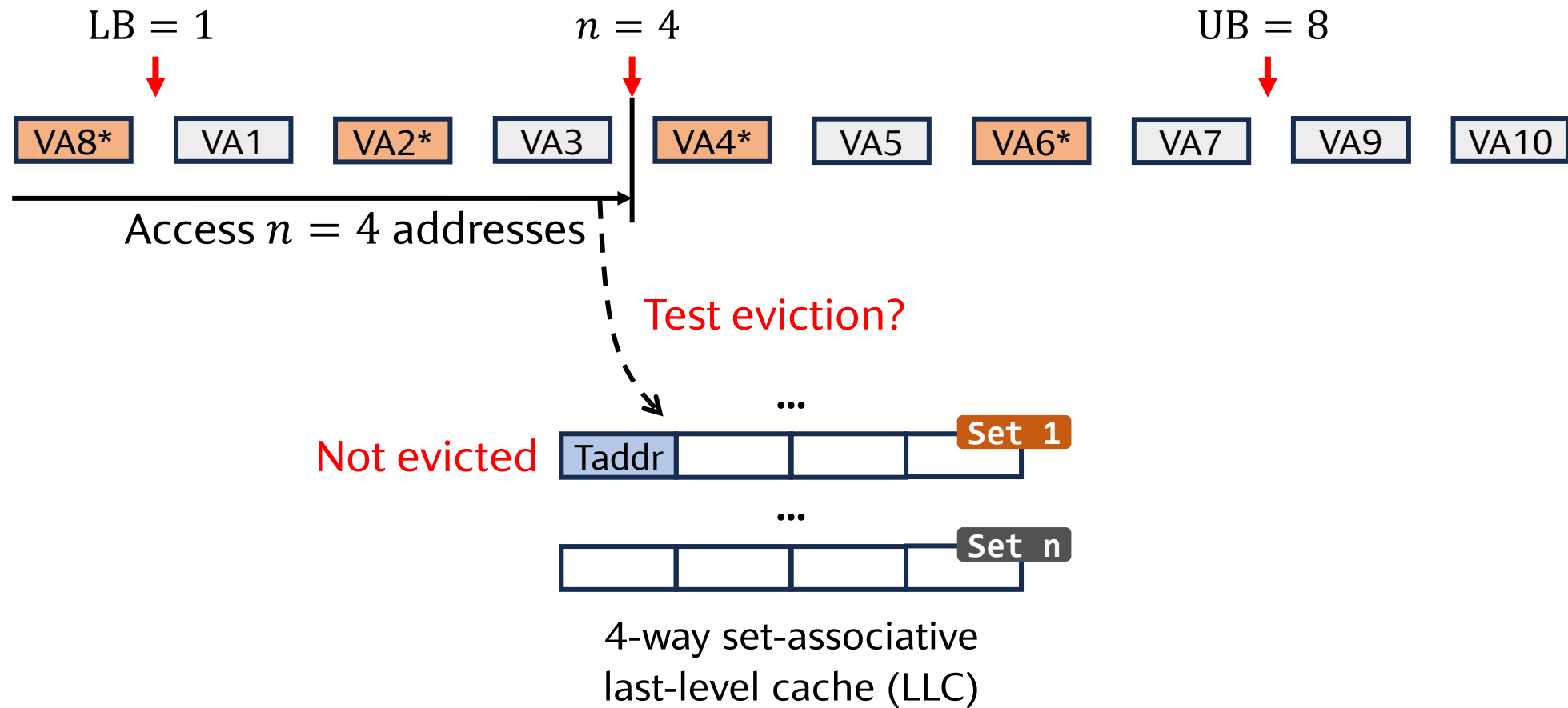


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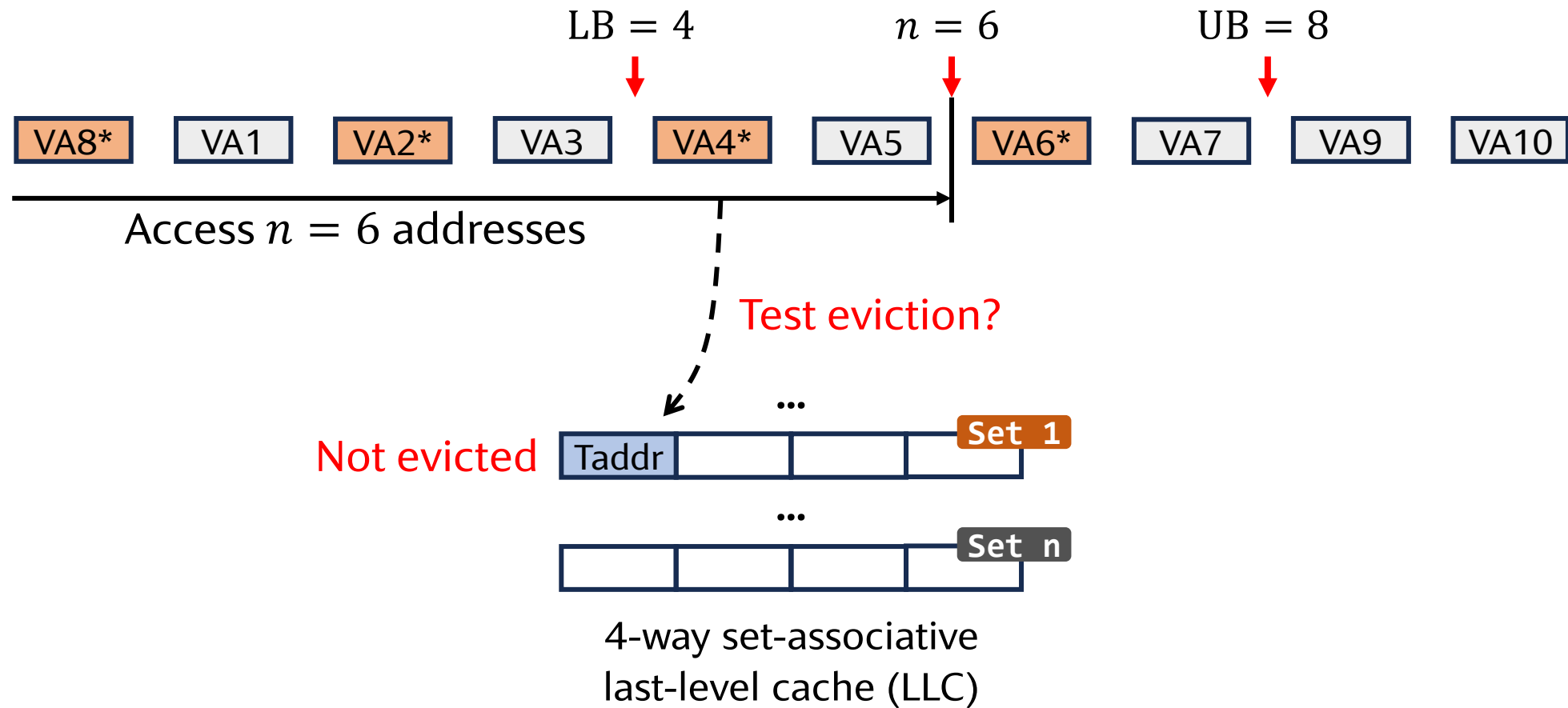


4-way set-associative
last-level cache (LLC)

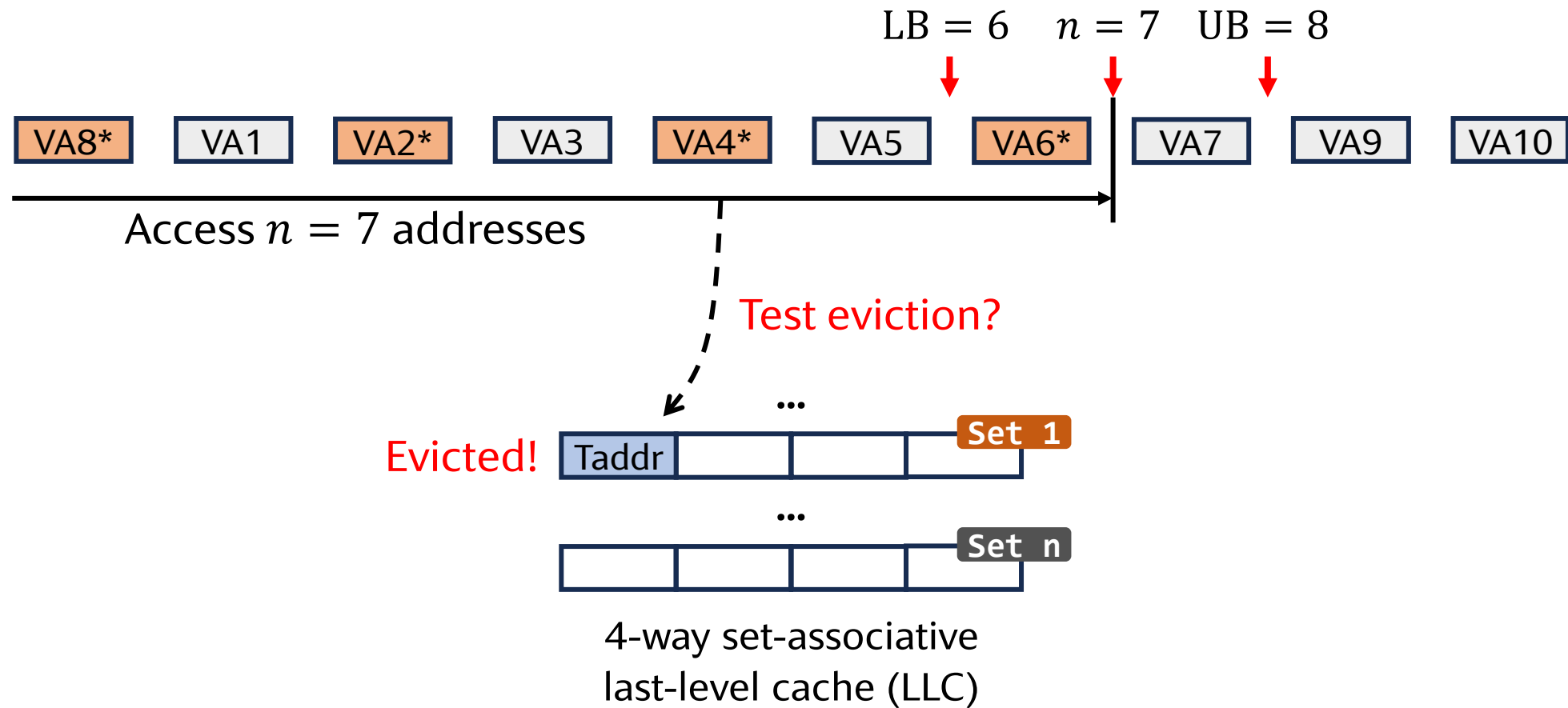
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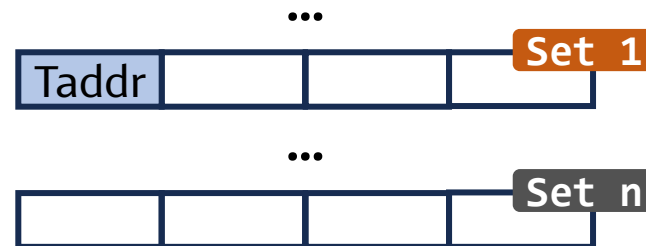
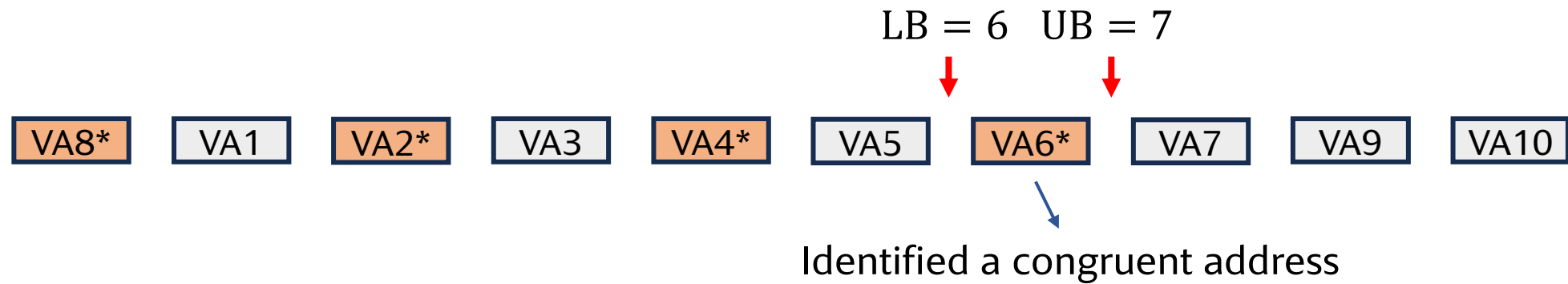
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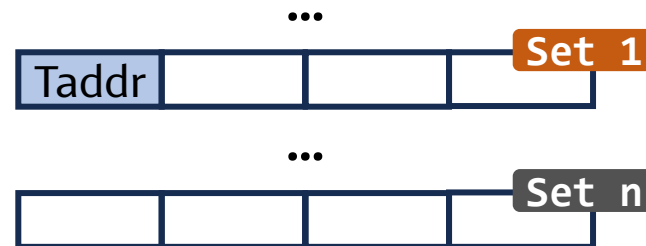
4-way set-associative
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New Technique 1: A Binary Search-Based Algorithm



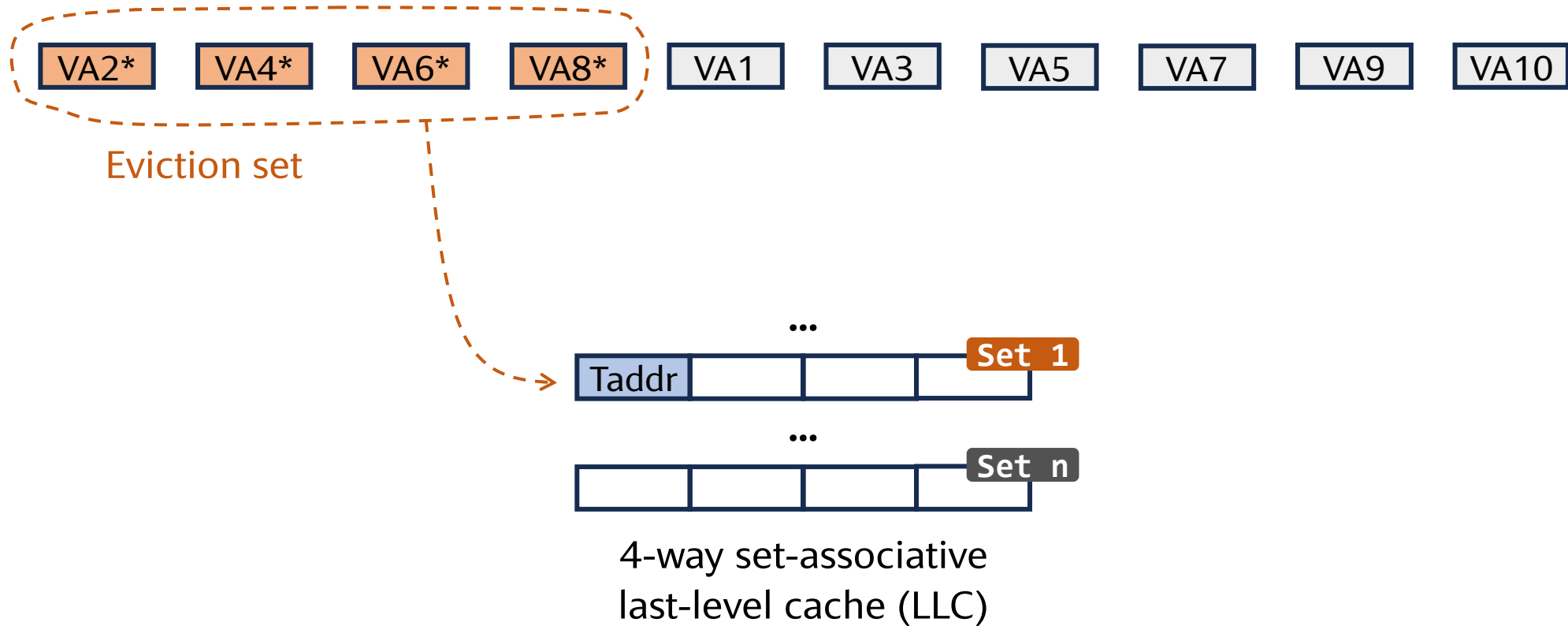
4-way set-associative
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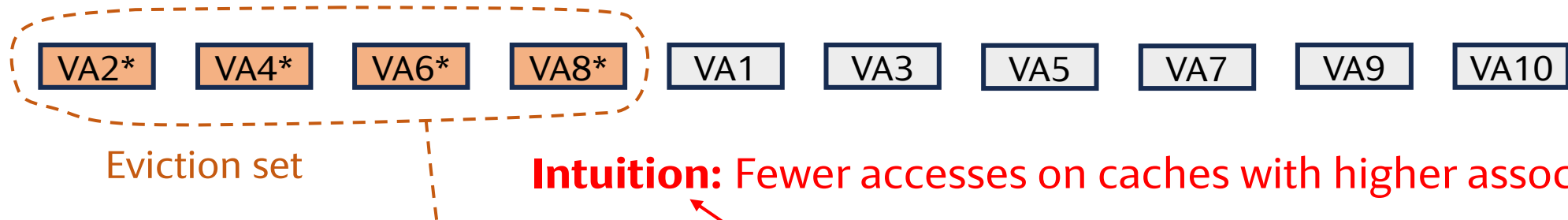


4-way set-associative
last-level cache (LLC)

New Technique 1: A Binary Search-Based Algorithm



New Technique 1: A Binary Search-Based Algorithm



Intuition: Fewer accesses on caches with higher associativity

Our algorithm requires $O(WN \log N)$ memory accesses

Group Testing requires $O(W^2N)$ memory accesses



4-way set-associative
last-level cache (LLC)

Intuition: Check out our paper for another optimization that reduces N

Attack Roadmap

- ✓ Step 1: Co-Locate / Zhao et al., Everywhere All at Once... (ASPLOS '24)
 - ✓ Step 2.1: Build Many Eviction Sets
 - ➔ **Step 2.2: Identify Target Set**
- } Challenges: Noise and dynamism

Insight: Victim accesses are periodic \Rightarrow Spectral analysis



↙
End-to-end, cross-tenant information leakage
in production Google Cloud

Attack Roadmap

- ✓ Step 1: Co-Locate / Zhao et al., Everywhere All at Once...
(ASPLOS '24)
 - ✓ Step 2.1: Build Many Eviction Sets
 - ✓ Step 2.2: Identify Target Set
 - ➔ **Step 2.3: Extract Information**
- } Challenges: Noise and dynamism

Insight: Overlap memory accesses of Prime/Probe to exploit memory-level parallelism
⇒ Low Prime/Probe latency ⇒ Good time resolution and noise resilience

Attack Roadmap

- ✓ Step 1: Co-Locate  *Zhao et al., Everywhere All at Once... (ASPLOS '24)*
- ✓ Step 2.1: Build Many Eviction Sets 

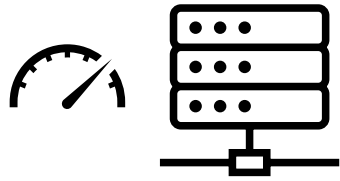
Target victim: A vulnerable ECDSA program from OpenSSL 1.0.1e

Result: Can extract an average of 68% of secret nonce bits (or a median value of 81%)

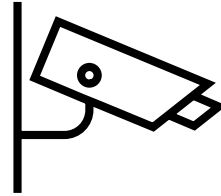
 End-to-end, cross-tenant information leakage in production Google Cloud

Conclusions

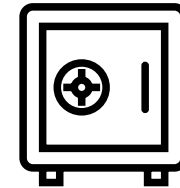
Cross-tenant information leakage with LLC Prime+Probe  Google Cloud



1. Fast LLC Channel Setup
> 10 hours → 2.4 minutes



2. Noise-Resilient Victim
Monitoring



3. Information Extraction

- + Google filed a **critical-level bug** to their product team
- + AWS revised their whitepaper on February 15, 2024

GitHub Repo: <https://github.com/zzrcxb/LLCFeasible>